Active learning of timed automata with unobservable resets

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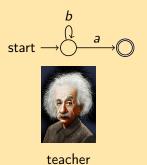
teacher



¹Angluin, "Learning Regular Sets from Queries and Counterexamples", 1987.

²Angluin, "Queries and Concept Learning", 1987.

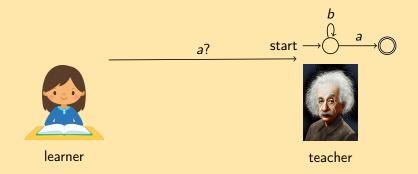






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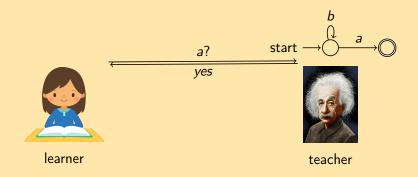
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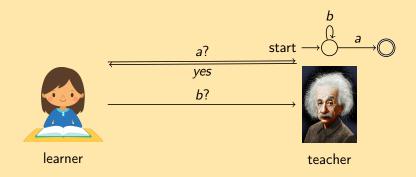
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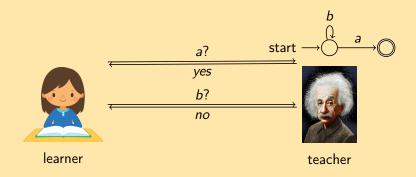
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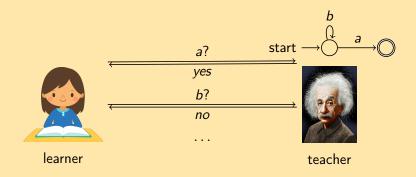
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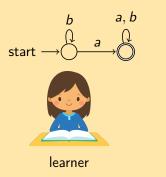
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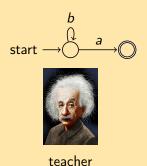




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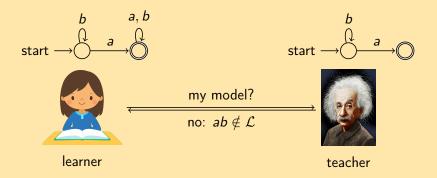






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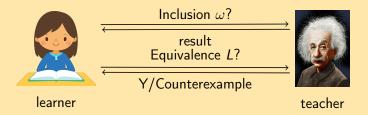
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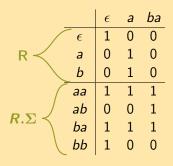


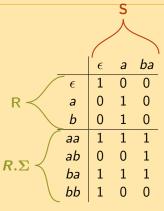


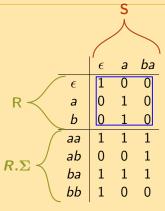
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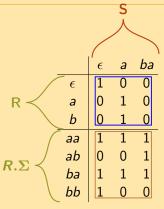
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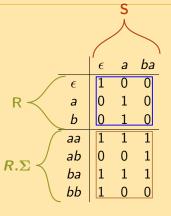
	ϵ	а	ba
ϵ	1	0	0
а	0	1	0
b	0	1	0
aa	1	1	1
ab	0	0	1
ba	1	1	1
bb	1	0	0



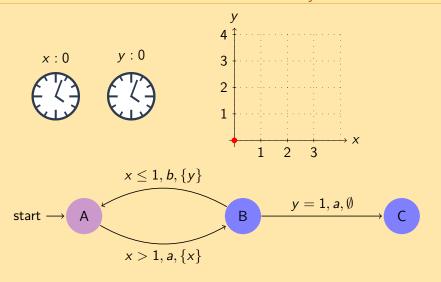




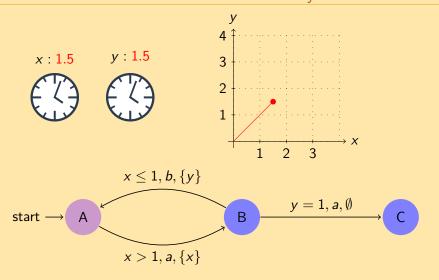




consistent
$$\forall u, v \in R, (u \sim_O v \Rightarrow \forall a \in \Sigma, u.a \sim_O v.a)$$

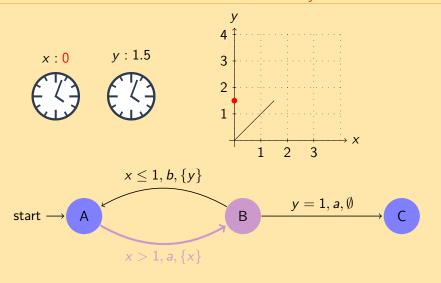


³Alur and Dill, "A Theory of Timed Automata", 1994.

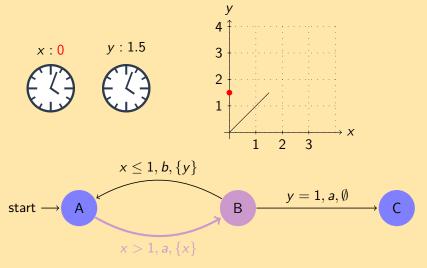




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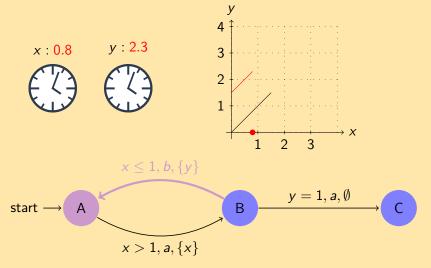


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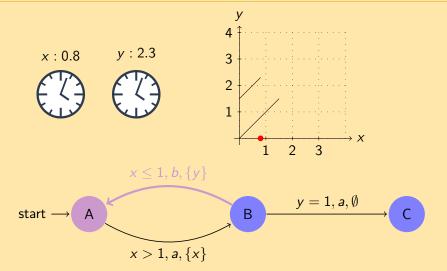
Execution: 1.5;a

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Execution: 1.5;a;**0.8;b**

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Execution: 1.5;a;**0.8;b**

⇒ Resets are unobservable



³Alur and Dill, "A Theory of Timed Automata", 1994.

Making resets observable

- ▶ One clock per letter in the alphabet.
- ► Each transition resets the clock corresponding to its letter.

⁴Alur, Fix, and Henzinger, "Event-Clock Automata: A Determinizable Class of Timed Automata",

Making resets observable

- One clock per letter in the alphabet.
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Observation: 0.7 a

1.4 *c*

8.0

⁴ Alur, Fix, and Henzinger, "Event-Clock Automata: A Determinizable Class of Timed Automata",

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Observation: $0.7 a [x_a \leftarrow 0] 1.4 c [x_c \leftarrow 0] 0.8$



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Dynamics:

$$\begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix} \xrightarrow{0.7} \begin{pmatrix} 0.7 \\ 0.7 \\ 0.7 \end{pmatrix} \xrightarrow{a} \begin{pmatrix} 0.0 \\ 0.7 \\ 0.7 \end{pmatrix} \xrightarrow{1.4} \begin{pmatrix} 1.4 \\ 2.1 \\ 2.1 \end{pmatrix} \xrightarrow{c} \begin{pmatrix} 1.4 \\ 2.1 \\ 0.0 \end{pmatrix} \xrightarrow{0.8} \begin{pmatrix} 2.2 \\ 2.9 \\ 0.8 \end{pmatrix}$$



⁴Alur, Fix, and Henzinger, "Event-Clock Automata: A Determinizable Class of Timed Automata",

Making resets observable

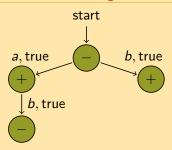
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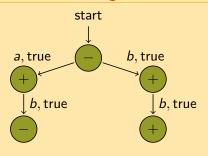






 $^{^{5}}$ Grinchtein, Jonsson, and Pettersson, "Inference of Event-Recording Automata Using Timed Decision Trees", 2006.

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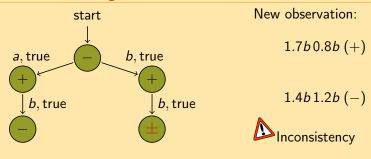
New observation:

1.7b0.8b(+)



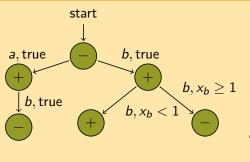
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New observation:

 $1.7b\,0.8b\,(+)$

1.4b1.2b(-)

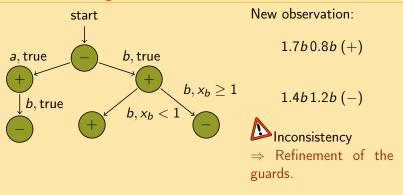
Inconsistency

⇒ Refinement of the guards.



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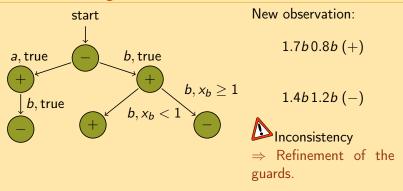


A consistent stucture can be folded as an ERA.



⁵Grinchtein, Jonsson, and Pettersson, "Inference of Event-Recording Automata Using Timed Decision Trees", 2006.

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A consistent stucture can be folded as an ERA.

- Interesting structure with good algorithms,
- Does not deal with reset guesses.



⁵Grinchtein, Jonsson, and Pettersson, "Inference of Event-Recording Automata Using Timed Decision Trees", 2006.

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Reset-free ERA

Making resets observable

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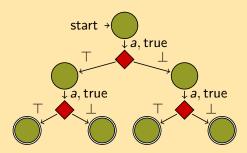


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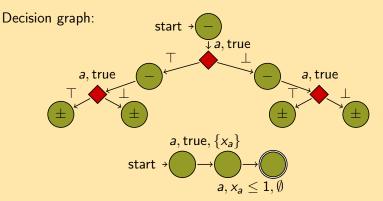
Observation: 0.7 a? 1.4 c? 0.8





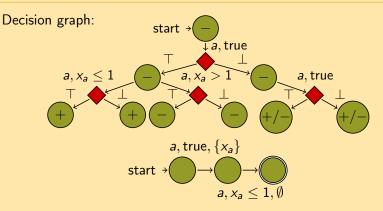
Active learning of RERAs: first try

Changing the structure



Active learning of RERAs: first try

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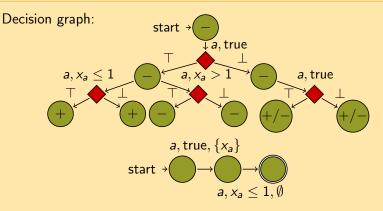


Inconsistency is detected ⇒ Separating guards



Active learning of RERAs: first try

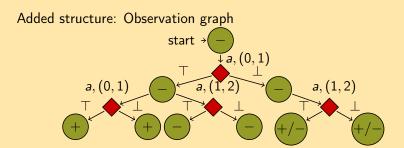
Changing the structure



Inconsistency is detected \Rightarrow Separating guards No separating guard without a reset: $0.9a\,0.8a\,(+)$ and $0.7a\,1.1a\,(-)$

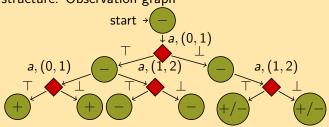


Guessing resets



Guessing resets

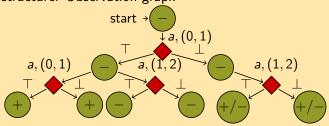
Added structure: Observation graph



► Each guard is "minimal": a cube of size one.

Guessing resets

Added structure: Observation graph

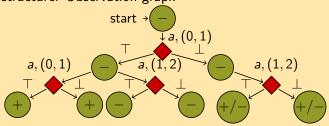


- ► Each guard is "minimal": a cube of size one.
- ▶ A node corresponding to both positive and negative observations is an **invalidity**.



Guessing resets

Added structure: Observation graph

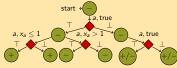


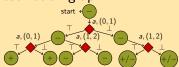
- ► Each guard is "minimal": a cube of size one.
- ▶ A node corresponding to both positive and negative observations is an **invalidity**.
- Allows to prune the main structure.



Combining the structures

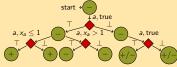
Decision graph:



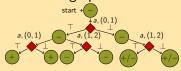


Combining the structures

Decision graph:



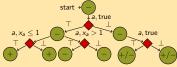
- Partitions the clock values
- Detects inconsistencies
- ► Infers guards



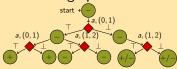
- Precise
- Detects invalidity
- ► Infers clock resets

Combining the structures

Decision graph:



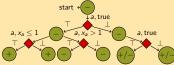
- Partitions the clock values
- Detects inconsistencies
- ► Infers guards
- ▶ Observations are propagated in both structures



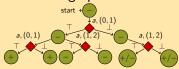
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Combining the structures

Decision graph:



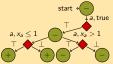
- Partitions the clock values
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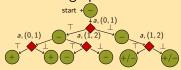
- Precise
- Detects invalidity
- ► Infers clock resets
- Observations are propagated in both structures
- ▶ Observation graph is used to prune the Decision graph

Combining the structures

Decision graph:



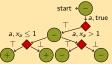
- Partitions the clock values
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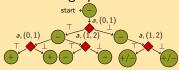
- Precise
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Combining the structures

Decision graph:



- Partitions the clock values
- Detects inconsistencies
- ► Infers guards



- Precise
- Detects invalidity
- ► Infers clock resets
- Observations are propagated in both structures
- Observation graph is used to prune the Decision graph
- Decision graph will be folded into a RERA

Conclusion

Learning of RERAs

- extend the learnable model classes;
- introduces the key notion of invalidity;
- paves the way to the learning of deterministic TAs.

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It requires

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It requires

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Interesting open questions:

- reduce the space-cost using implicit structures;
- ▶ redact and analyse the generalization to deterministic TAs.

