ALPHA:
A functional data-parallel language based on Polyhedra

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Outline

- Introduction
- ALPHA Syntax and Semantics (of Domains)
- Substitution & Normalization
- Change of Basis
- Conclusion

The Essence of Data Parallelism

- Collectionsof elementary data objects
- Pointwise Operations
- Alignmentof different collections
- Conditional operations on (parts of) collections
- Reductions/Scansof associative and commutative operators

Alpha Syntax

```
system   ⟨Name⟩:     ⟨Domain⟩
                     ( [ ⟨Var⟩* : ⟨Domain⟩ of ⟨Type⟩ ]* )
           returns   ( [ ⟨Var⟩* : ⟨Domain⟩ of ⟨Type⟩ ]* );
var
                     [ ⟨Var⟩* : ⟨Domain⟩ of ⟨Type⟩ ]*

let
      [ ⟨Var⟩= ⟨Expression⟩]*
tel
```

Polyhedral Domains (Dual Representation)

- Integral Polyhedron: All points that **satisfy** a finite number of linear inequalities

$$\{z \in \mathcal{Z}^n \mid Qz \geq q\}$$

- Generators: All **convex combinations** of a set of vertices

$$\{z \in \mathcal{Z}^n \mid z = a^T G;\ \sum_i a_i = 1\}$$

- Alpha Syntax: { ⟨Idx⟩* | [⟨IdExpr⟩ >= ⟨IdExpr⟩]* }

Each Representation is useful in certain contexts. A standard (self dual) algorithm to convert from one to other.

Syntax of Expressions

- ⟨Var⟩| ⟨Const⟩ — Atomic
- ⟨Expression⟩op⟨Expression⟩ — Pointwise Ops
- ⟨Domain⟩ : ⟨Expression⟩ — Restriction
- case [⟨Expression⟩]* esac — Case
- ⟨Expression⟩. ⟨Dep⟩ — Affine Dependency

  where ⟨Dep⟩ has syntax: (⟨Idx⟩* -> ⟨IdExpr⟩*)

- reduce (op, ⟨Dep⟩ , ⟨Expression⟩) — Reduction

Example (Fibonacci)

```
system FibSys :  {N| 1<=N} () returns(F : integer );
var    Fib :     {i | 1 <= i <= N} of integer
let
 Fib =
       case
       {i| i<=2} :  1.(i -> ) ;
       {i| i>=3} :  Fib.(i -> i-1) + Fib.(i -> i-2) ;
       esac ;
 F =   Fib.( ->N) ;
tel ;
```

Semantics of Alpha Expressions

- Each expression denotes a function from a domain, D to a ⟨Type⟩ (real, integer or boolean)
- The Semantic function is compositional, and straightforward (eg, `A + B` denotes the sum of `A` and `B`) . . .
- Domains too are straightforward . . .
- except for affine dependency

Domains of ALPHA Expressions

- Constantsand Variables $\boxed{Dom(c) = \mathcal{Z}^0}$ $\boxed{Dom(V)}$
- Pointwise operators $\boxed{Dom(e_1 \oplus e_2) = Dom(e_1) \cap Dom(e_2)}$
- Restriction $\boxed{Dom(D : e) = D \cap Dom(e)}$
- **case** operator $\boxed{Dom(\textbf{case } e_1; ..; e_n; \textbf{esac}) = \bigcup_{i=1}^{n} Dom(e_i)}$

Affine Depencency

- Q: What does the expression $e.(z \rightarrow f(z))$mean?
- A: An expression whose value at $z$is the same as the value of $e$at $f(z)$

  Example: `Y = X.(i,j -> i)` or `Y[i,j] = X[i]`

- Its domain is $\boxed{Dom(e.f) = f^{-1}(Dom(e))}$

Substitution

```
substituteInDef[Fib, Fib, 1]; show[]
```

```
system FibSys :  {N| 1<=N} () returns(F : integer );
var    Fib :     {i | 1 <= i <= N} of integer
let
 Fib =
       case
       {i| i<=2} :  1.(i -> ) ;
       {i| i>=3} :    case
                      {i | i<=2} :  1.(i->);
                      {i | 3<=i} :  Fib.(i->i-1) + Fib.(i->i-2);
                      esac.(i -> i-1) + Fib.(i -> i-2) ;
       esac ;
 F =   Fib.( ->N) ;
tel ;
```

Normalization: syntax of a normal expression

- ⟨BasicExpr⟩::= ⟨Var⟩| ⟨Const⟩| ⟨Var⟩.⟨Dep⟩ | ⟨Const⟩.⟨Dep⟩
- ⟨InnerExpr⟩::= ⟨BasicExpr⟩| ⟨BasicExpr⟩op ⟨InnerExpr⟩
- ⟨NormalExpr⟩::=
  - ⟨InnerExpr⟩| ⟨Domain⟩ :⟨InnerExpr⟩|
  - case [⟨InnerExpr⟩]* esac
  - case [⟨Domain⟩ :⟨InnerExpr⟩]* esac

Normalization Rules (subset)

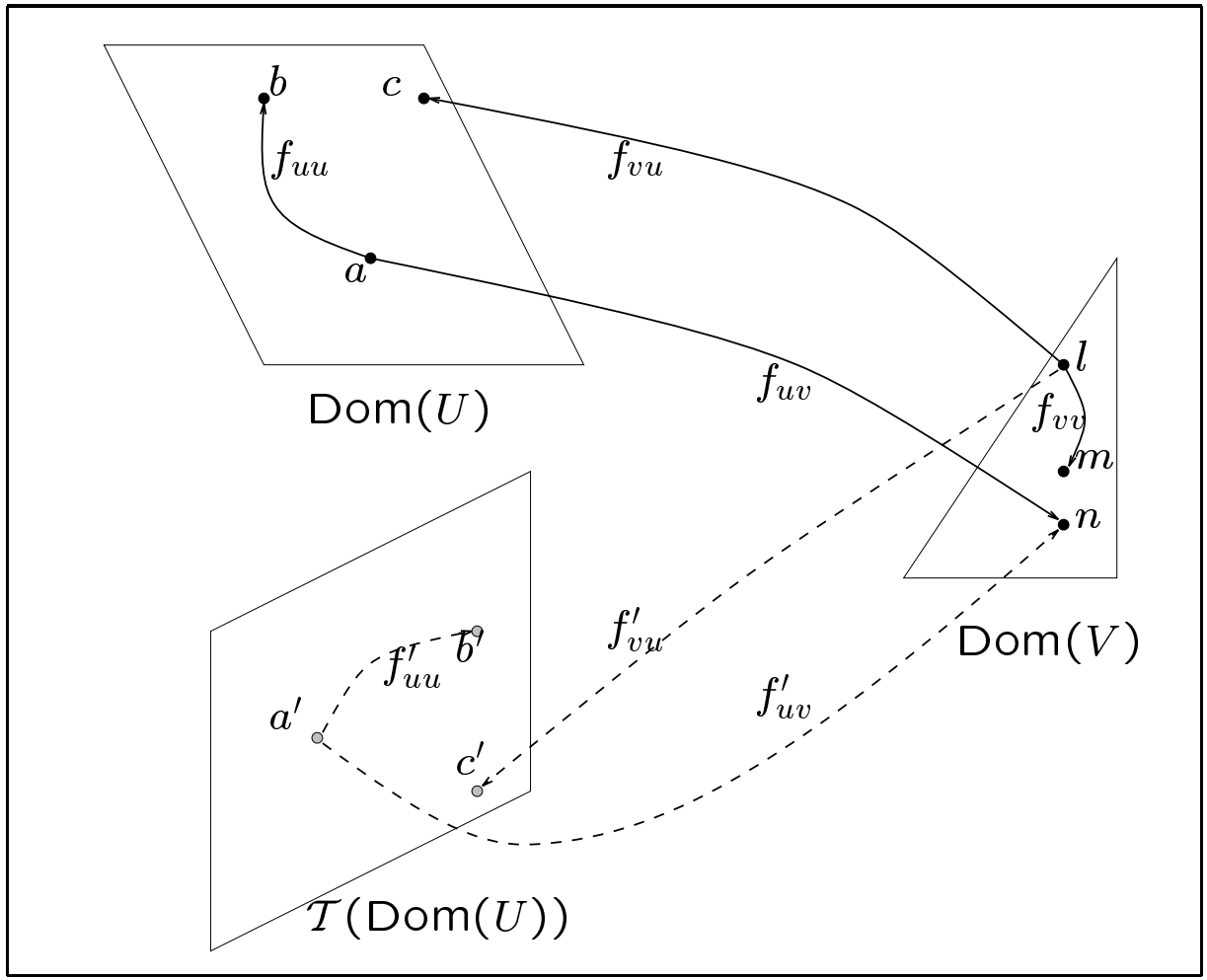
No.	Replace	By	Comments
1	$e.f$	e	if $f(z) = z$
2	$(e_1 \oplus e_2).f$	$(e_1.f) \oplus (e_2.f)$	
3	$e \oplus$ **case** e_1; $\vdots$ e_n; **esac**	**case** $e \oplus e_1$; $\vdots$ $e \oplus e_n$; **esac**	
4	$(D : e_1) \oplus e_2$	$D : (e_1 \oplus e_2)$	
5	$e_1 \oplus (D : e_2)$	$D : (e_1 \oplus e_2)$	
6	$(e.f_1).f_2$	$e.f$	$f = f_1 \circ f_2$
7	$D_1 : (D_2 : e)$	$D : E$	$D = D_1 \cap D_2$
8	$(D : e).f$	$D' : (e.f)$	$D' = f^{-1}(D)$

Normalization, example

```
normalize[]; show[]
```

```
system FibSys :  {N| 1<=N} () returns(F : integer );
var    Fib :     {i | 1 <= i <= N} of integer
let
 Fib =
       case
       {i| i<=2} :  1.(i -> ) ;
       {i| i=3}  :  1.(i ->) + Fib.(i -> i-2) ;
       {i| i>=4} :  Fib.(i -> i-2) + Fib.(i -> i-3)
                        + Fib.(i -> i-2) ;
       esac ;
 F =   Fib.( ->N) ;
tel ;
```

Change of Basis



Change of Basis Transformation Rules

- Space time transformations – applied to any (local) variable, V in an ALPHA system
- Unimodular transformation, T (specified like a dependency function)
- Three simple rewrite rules: Replace
  - the domain, D of V by $T(D)$
  - all occurrences of V on the rhs of all equations by $V.T$
  - E, the *entire rhs* of the equation defining V by $E.T^{-1}$

Example (Fibonacci again)

```
load["fib.a"]; changeOfBasis["Fib.(i -> -i)", "i"]; show[]
```

```
system FibSys :  {N| 1<=N} () returns(F : integer );
var    Fib :     {i | -N<=i<=-1}of integer
let
 Fib =
       case
       {i| i<=2} :  1.(i -> ) ;
       {i| i>=3} :  Fib.(i->-i).(i -> i-1) + Fib.(i->-i).(i -> i-2) ;
       esac . (i->-i);
 F =   Fib . (i->-i). ( ->N) ;
tel ;
```

Normalize again

```
normalize[]; show[]
```

```
system FibSys :  {N| 1<=N} () returns(F : integer );
var    Fib :     {i | -N<=i<=-1}of integer
let
 Fib =
       case
       {i| -2<=i}:  1.(i -> ) ;
       {i| i<=-3}:  Fib.(i -> i+1)+ Fib.(i -> i+2);
       esac ;
 F =   Fib.(->-N);
tel ;
```

Why does it work?

ALPHA domains (unions of finitely many polyhedra) constitute an abstract data type (ADT), closed under:

- Intersection
- Union (of finitely many members of the ADT)
- Preimage by (arbitrary) affine functions (the class of dependencies in ALPHA).
- Image by unimodular affine functions (the class of transformations used for changes of bases).