

Classes of Rational Graphs

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Irisa

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Graphs

Graphs

A **graphe** G is a subset of $V \times A \times V$

- V is an arbitrary set of **vertices**
- A is a finite set of **labels**

Conventions

(u, a, v) is denoted by $u \xrightarrow{a} v$

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For rational graphs $V = \{0, 1\}^*$ (i.e. some finite alphabet)

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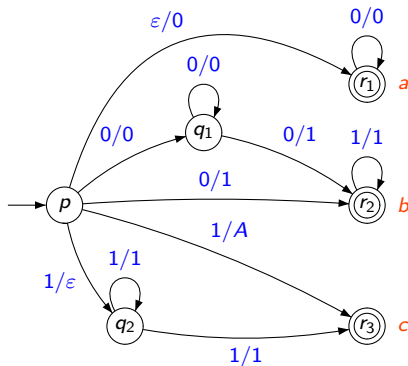
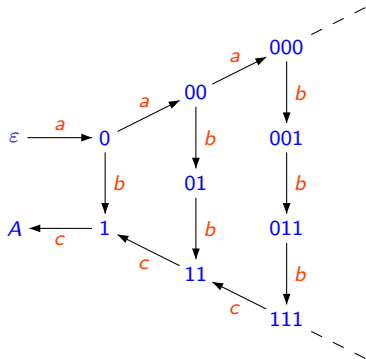
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Rational graphs

Definition

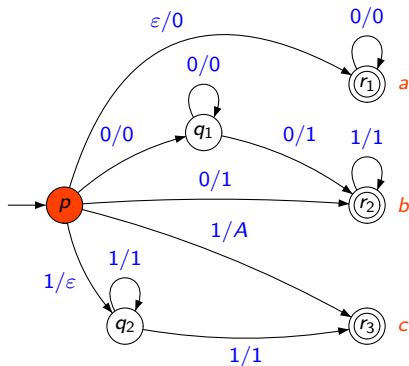
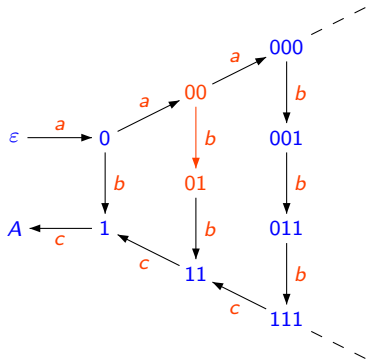
A rational graph is defined by a rational transducer



Rational graphs

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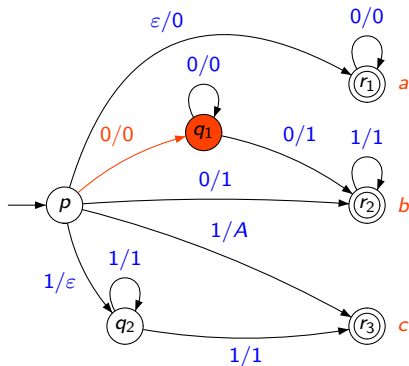
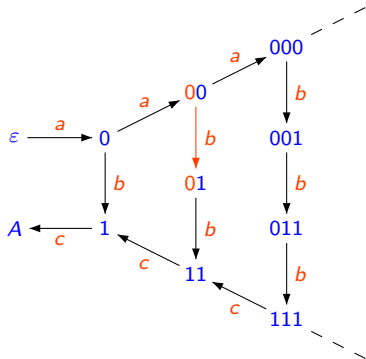
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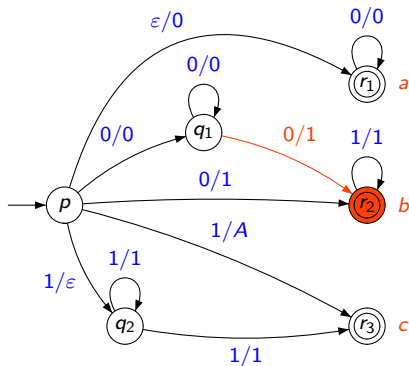
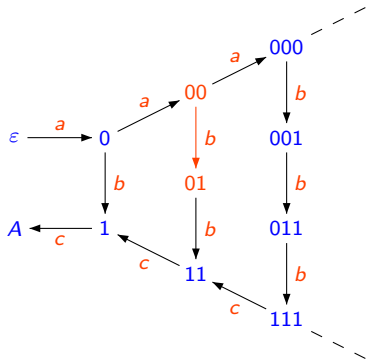
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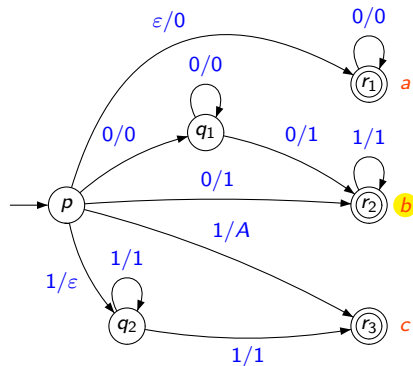
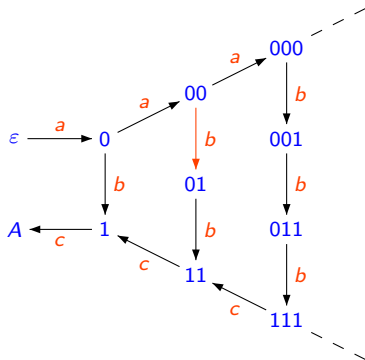
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Rational graphs

Definition

A rational graph is defined by a rational transducer



Classical questions

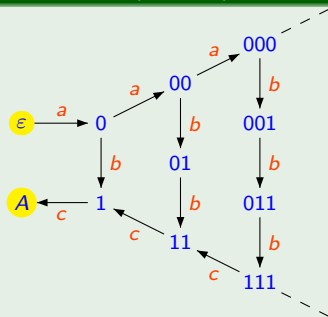
Proposition

Following properties are *undecidable* for rational graphs:

- (i) *Accessibility (between two vertices);*
- (ii) *Connectedness (of the whole graph);*
- (iii) *Isomorphism (of two graphs);*
- (iv) *First order theory (of a rational graph).*

Traces

Example (traces)



The trace of this graph between ε and A is $\{a^n b^n c^n \mid n \in \mathbb{N}\}$

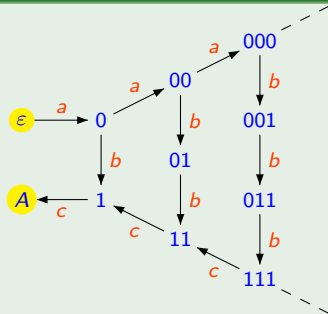
$$\text{Tr}(G, \varepsilon, A) = \{a^n b^n c^n \mid n \in \mathbb{N}\}$$

Theorem (Morvan Stirling 01)

The traces of rational graphs between two vertices coincide with the context sensitive languages.

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Deterministic Rational Graphs

Definition

A rational graph is **deterministic** if each vertex is the source of at most one arc for each label

Proposition

*Determinism is **decidable** for rational graphs*

Proposition

*First order theory and accessibility are **undecidable** for deterministic rational graphs*

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Traces

Proposition

*The **traces** of deterministic rational graphs form a boolean algebra of deterministic context-sensitive languages containing non context-free languages.*

Question

Are there context-free languages not contained in the traces of deterministic rational graphs?

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Are there **context-free** languages not contained in the traces of deterministic rational graphs?

Rational Trees

Definition (Tree)

- at most one ancestor per vertex
- a single root
- connected

decidable

decidable

undecidable

Rational Trees

Definition (Tree)

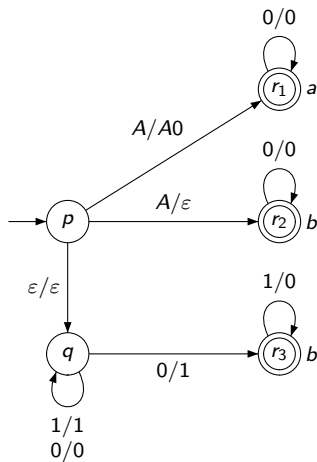
- at most one ancestor per vertex decidable
- a **single** root decidable
- connected undecidable

Rational Trees

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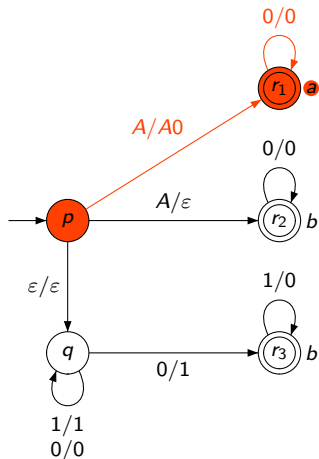
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A simple example - 2^n -tree

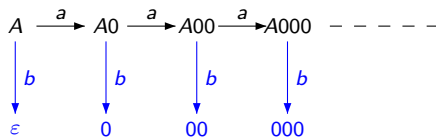
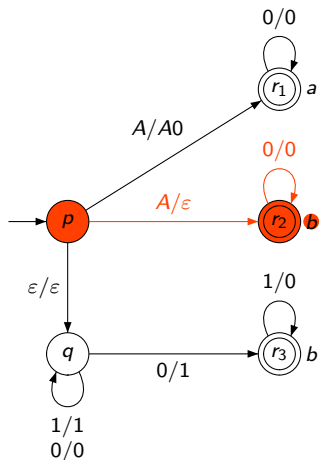


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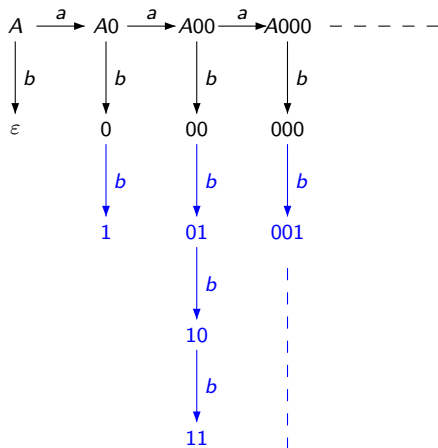
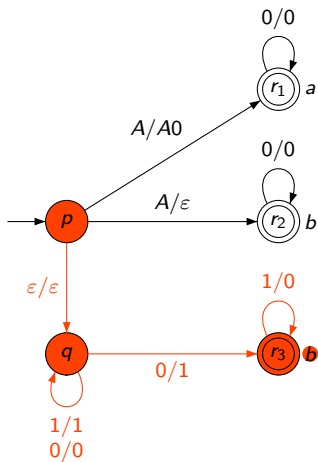
$$A \xrightarrow{a} A0 \xrightarrow{a} A00 \xrightarrow{a} A000 \dots$$



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Properties of rational trees

Proposition

For rational trees

- *inclusion and equality is **decidable***
- *accessibility (between vertices) is **decidable***
- *effective rational set of leaves*
- *some **sub-trees** are non-rational*

Theorem (Carayol-Morvan '06)

Properties of rational trees

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Theorem (Carayol Morvan 06)

- *First order theory of rational trees decidable*

→ First order theory of rational rooted ordered trees

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Traces

Definition

The **trace** of a tree is the set of path labels from the **root** to the **leaves**

Conjecture

There are context-free languages which are not the trace of any rational tree

Example

The set of words having the same number of a 's and b 's seems not to be the trace of any rational tree

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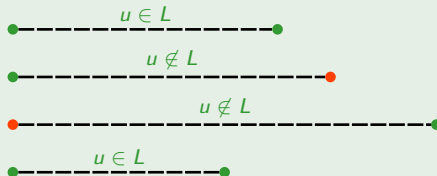
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Curious fact

Example

Given any context-sensitive language L we may construct such a rational graph (green and red states form rational sets)



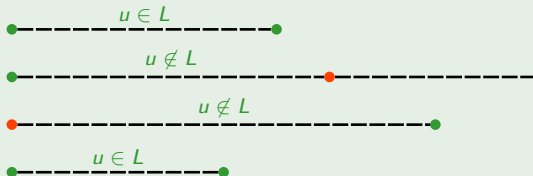
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Given any context-sensitive language L (with alphabet X), there is a rational tree T labelled on $X \cup \{\#\}$ such that $\text{Tr}(T) \cap X^ = L$*

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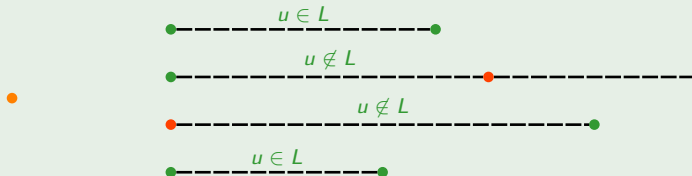
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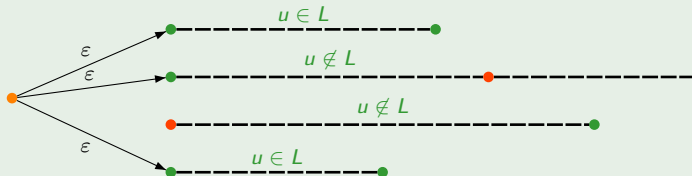
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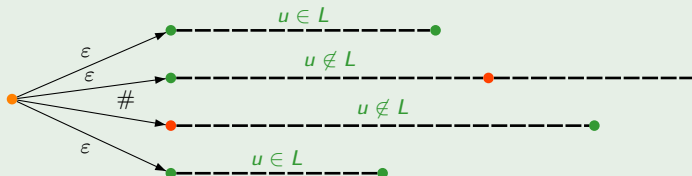
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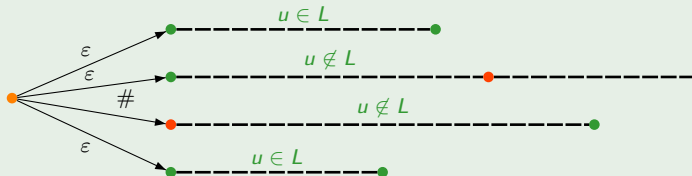
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Rational forests

Proposition

*It is **undecidable** to know if a rational graph is a forest*

Proposition

*The first order theory of rational forests is **decidable***

Proposition

The trace of rational forests between a root and leaves coincide with context-sensitive languages

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Rational Directed Acyclic Graphs

Definition

A rational DAG is a rational graph with no cycle

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It is *undecidable* to know if a rational graph is a DAG

Proposition

There is a connected rational DAG with undecidable first-order theory

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Automatic Graphs

Definition

A graph over $X^* \times A \times X^*$ is automatic if it is recognized by a letter-to-letter transducer followed by a terminal function taking values in

$$(\text{Rat}(X^*) \times \{\varepsilon\}) \cup (\{\varepsilon\} \times \text{Rat}(X^*))$$

Example

The previous examples ($a^n b^n c^n$ -graph and 2^n -tree) are automatic

Proposition (Hodgson 83)

The first order theory of automatic graphs is decidable

Theorem (Rispal 01)

The traces of automatic graphs are context-sensitive languages

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Monotonous Rational Graphs

Definition

A rational graph is **monotonous** if for each transition u/v of its transducer either satisfy $|u| \leq |v|$ or $|u| \geq |v|$

Proposition

Accessibility is decidable for monotonous graphs

Question

Is the first order theory of monotonous graphs decidable?

Monotonous Rational Graphs

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Traces of monotonous graphs

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From each rational graph a bisimilar monotonous graph is computable

Corollary

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Commutative Rational Graphs

Definition (Commutative graphs)

A rational graph is **commutative** if its transducer is defined over a commutative monoid $(\mathbb{N}^n, \mathbb{Z}^n)$

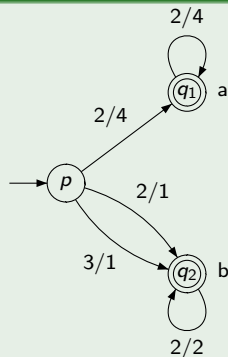
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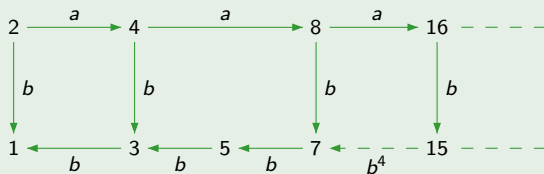
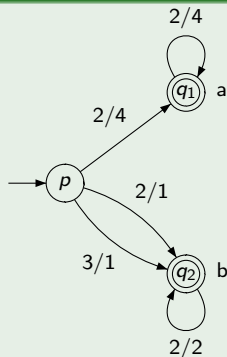


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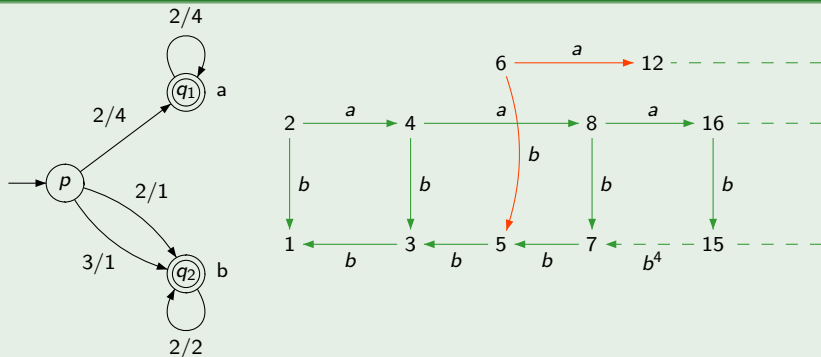
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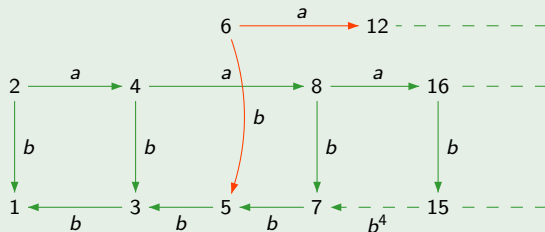
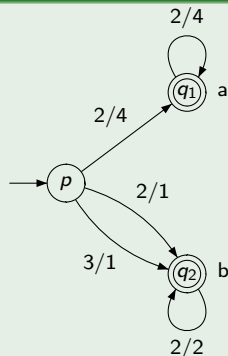


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Example (A commutative rational graph)



$$Tr(G, 2, 1) = \{a^n b^{2^n} \mid n \in \mathbb{N}\} \notin CF(\{a, b\}^*)$$

Some results

Theorem (Eilenberg, Schützenberger 69)

The family of rational sets in a commutative monoid form a Boolean algebra

Corollary

*First order theory is **decidable** for commutative graphs*

Proposition

*Accessibility is **undecidable** for commutative rational graphs*

Proposition

The families of commutative rational graphs and prefix recognizable graphs are incomparable

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Conclusion

Conclusion

Several families of rational graphs

- Structural

Deterministic rational graphs

Rational trees

Rational DAG

- Automata driven

Future work

- Some families are promising

Deterministic rational graphs, Monotonous rational graphs,

Commutative rational graphs

Combinatorial such as family \mathcal{R} , \mathcal{R}^* , \mathcal{R}^{\dagger}

Conclusion

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Several families of rational graphs

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Deterministic rational graphs, Monotonous rational graphs,
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- Exhibit a such a family with FO+ACC decidable

Conclusion

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Several families of rational graphs

- Structural

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Rational DAG

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